ECE 447 Fall 2025

Lesson 37
Error Correction:
Parity, CRC, and
Hamming Codes



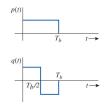
SCHEDULE AND ADMIN

- Schedule.
 - Lesson 37 Error correction
 - Lesson 38 MATLAB Lab 7: Matched filters, multi-path, OFDM, BER (workday - no attendance for class)
 - Lesson 39 Advanced topics: OFDM, MIMO, CDMA
 - Lesson 40 Course review
- Admin
 - HW8. Due 02 Dec (Lsn 39) to Gradescope.

REVIEW

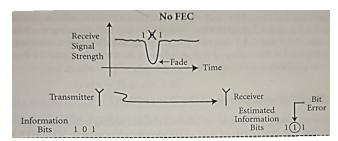
Schedule and Admin

 Example problem: Find the BER for a system using polar **signaling** pulses $\pm 3p(t)$ with $T_b = 1\mu s$ and $N_0 = \frac{25\mu W}{H_Z}$.



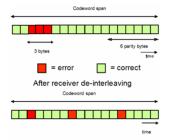
ERROR CORRECTION

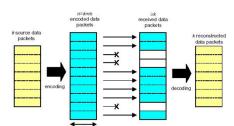
- Bit errors occur for a variety of reasons, but can correct many by pre-processing bits at the transmitter forward error correction (FEC)
- "Error correction coding requires a strong mathematical background" so true!
- Types of codes (in order of complexity): repetition, block, convolutional, Viterbi algorithm, turbo, low-density parity check (LDPC), polar



- First, though errors often consecutive in time and/or frequency
- Interleaving spreads out code symbol errors at receiver decoder receives few to zero consecutive errors
- Synchronization required!

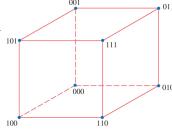
Review





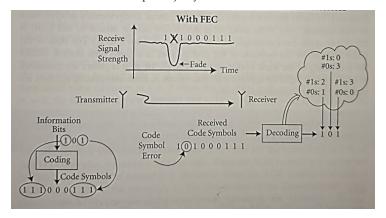
ERROR CORRECTION BASICS AND TERMINOLOGY

- Code rate: ratio of data bits k to total bits n creates an (n, k) code
- *d* is bit vector **message** (data)
- *c* is bit vector **codeword**
- Number of redundant (aka check aka parity) bits is m = n k
- Hamming distance: number of bits that are different between two bit vectors
- Hamming space: multi-dimensional cube; spheres drawn around possible codeword vertices - the radius spheres can have without overlapping equals the number of errors that can be corrected
- To *correct* t errors, min Hamming distance $d_{min} = 2t + 1$
- To *detect t* errors, min Hamming distance $d_{min} = t + 1$



REPETITION CODE

- Repeat each bit a certain number of times, then transmit
- To decode at receiver, use a simple majority rule



HAMMING CODE

- Type of linear block code
- c = dG, where G is called the **generator matrix**
- We transmit c and receive r, which may have errors
- $s = rH^T$ is called the **syndrome**, where *H* is the **parity check matrix**
- The row in H^T that matches s is the bit of the codeword in error flip that bit
- If s = 0, then no errors
- Coding gain: reduction in req'd SNR to achieve specific BER

WHAT DO WE "GAIN"?

• Coding gain: reduction in req'd SNR to achieve specific BER

